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Lecture 5: Proof of Nash's Theorem
Wednesday, 7 September 2022
                             9:45 AM
Theorem (Nash, 1950): Every game w/a finite set of players & strategies
  has an equilibrium in mix ed stralegies
  - proves existence, some properties, but not constructive
  - usus "fixed-pt. theorems"
Theorem (Brouwe's fixed- Pt. Theorem): Let XCR" be
     convex and compact. Let \varphi: X \to X be a continuous
      function. Then q has a fixed point in X, i.e., I ne X
      s.t. cp(n) = n.
   Intuitively, let \Gamma = (N, \{S_i\}_{i \in \mathbb{N}}, \{u_i\}_{i \in \mathbb{N}}) be an n-per son game.
   Suppose each playe has in pure stalique.
   Then for each player i, \Sigma_i = \Delta_m = \{n \in \mathbb{R}^m : \Sigma_{x_i} = 1\} \subseteq \mathbb{R}^m
            2 = 2, \times ... \times 2n \leq R^{m \times n}
    Consider a function p! 2 -> I , oblined as:

\varphi\left(\sigma_{1},\sigma_{2},\ldots,\sigma_{n}\right) = \left(\sigma_{1}',\sigma_{2}',\ldots,\sigma_{n}'\right)

    where, if o = (o, ..., on),
                       oi is the bust-rusponse to oi
                  i.e., \sigma_i' = arg max y_i(n, \sigma_i)

n \in \Sigma_i
    Then, if or is a fixed-pt. for p,
               i.e., \varphi(\delta_1, \delta_2, \ldots, \delta_n) = (\delta_1, \ldots, \delta_n)
     the or is a NE.
    (Problem: note that there can be infinite best-ousponses, so
                         q is not well-defined
      So let's do this more formally.
 ① X \subseteq \mathbb{R}^n is convex if, n_1, n_2 \in X \Rightarrow \lambda n_1 + (1-\lambda) n_2 \in X
       4 XE [0,1]

    X ⊆ R<sup>n</sup> is <u>bounded</u> if ∃ M ∈ R s.t. X ⊆ B(0, m)

                                                                                = \{ n \in \mathbb{R}^n : \| n \|_2 \leq M \}
(ii) X is closed if it contains all its limit points:
        (a) If n., n., ... is a converget sequence, n: EX,
                and \lim_{i \to \infty} x_i = \hat{x}, then \hat{x} \in X.
      X is open if In EX, JE>O s.t. B(n, E) = {y ER?: ||n-y||_2
                                                                                        < E } C X
         (b) X is closed if IRM X is open
 (i) f is continuous ((E.S)-untinity) if the EX, te SO, 75>0
         s.t. \forall y \in X : \|y - x\|_2 \leq \delta, \|f(y) - f(n)\|_2 \leq \varepsilon
 Examples & Counter examply
      X \leq Q:
         1. not bounded: X = {nell:n>0}
                not chosed: X = [o,1)
          3. not convex : X = [0,1] U[2,3]
    BFPT in 1-dimension
          If X is convex & compact, then X= [a,b]
          Let f be a continuous fr. If f(a) +a, f(b) + b,
                                                  then f(a) > a, f(b) < b
                                                   =) f(a)-a>0, f(b)-b<0
         Let g(n) = f(n) - \kappa, n \in [a,b]. The since
                g(a) > 0, g(b) < 0, by intermediate value theorem,
                 In e [a,b] s.t O= g(n) = f(n) - n. This is a fixed pt.
      Examples of no fixed-pt:
         - f continuous, X not bounded;
               X = [0, \infty), f(n) = 2n
          - f continuous, x not closed:
                    X = (0,1], f(x) = x(2)
          - f continuous, X not convex:
                      X = [0,1] \cup [2,3], f(n) = \begin{cases} n+2 & \text{if } n \leq 1 \end{cases}
                                                         \begin{cases} n-2 & \text{if } n \geqslant 2 \end{cases}
         - f' dis continuous, X compact & convex:
                    X = [0,1], f(n) = \begin{cases} n + 1/2 & n < 1/2 \\ n - 1/2 & n > 1/2 \end{cases}
Proof of Nash's Thorem via BFPT
 (Note: Can prove Nach's Theorem using other fixed-pt. theorems,
    eg. Kakutani fixed-Pt. Theorem, Will stick to BFPT for now)
    n-player game (N, {Si]iGN, {40}iGN)
     tor simplicity, assume each player has on pure strategies
    Recall: Zi = Dm CR , oi & Zi
                    2 = \sum_{i} x_{i} \times \sum_{n} x_{i} = \sum_{i} \sum_{n} R^{n \times n}
     Note: 2 is compact & convex (prove your self)
                  Given 6, 6-i, ui(6), ...
                   \Theta:(n,\sigma)=u:(n,\sigma_{-i})-\|n-\sigma_{c}\|_{2}^{2} (n\in\Sigma_{i})
    Defn:
                      " distance - adjusted whility"
                   \varphi_i(\varepsilon) = \underset{n \in \Sigma_i}{\text{avg max}} \left[ u_i(n, \varepsilon_i) - || n - \sigma_i ||_{n}^{2} \right]
                    i.e., Pi(0) is the mixed strategy for i that maximilys
                    the distance - aggisted outily.
                     (will show later that q_i(s) is well-defined)
              \varphi(6) = (\varphi_i(6), \varphi_2(6), \ldots, \varphi_n(6))
    Need to Show:
           1) Pi(6) is well-defined
          (1) P, 2 Satisfy conditions for RFPT
          (1) fixed-pt. for p is a NE for the game.
 For D, will show that "distance - adjusted utility" Di(n, 0) is
   Strictly Concave in a (keeping of fixed), and hence Di(1,0) has a
   Unique marcinizer (ande hence qi (6) is well-defined)
Lemma! HIEN, HOEZ, Di (n, o) is strictly concave in the
                      first argument.
       Since \theta_i(n, \sigma) = u_i(n, \sigma_i) - ||n - \sigma_i||_2, let's consider flese
Claim: The function ui (n, oi) is linear in the first argument
 Kroof: Let Vi (Sk, o=i) be the expected whilly for the 6th
                 pure strategy, given si These are fixed, given si
                  Then ui (n, si) = \( \sum_{\chi} \mathbb{n}_{\chi} \mathbb{n}_{\ch
                  Hence linear.
   Also, \| \mathbf{n} - \mathbf{\sigma}_i \|_2^2 = \sum_{k=1}^{\infty} (\mathbf{n}_k - \mathbf{\sigma}_i(k))^2 is shirtly conex
   Hence, \Theta_i(n, \sigma) = U_i(n, \sigma_i) - |n - \sigma_i|_{L^{\infty}} is strictly concave.
    Since a strictly concave for, on a convex set has a
   unique neximizer, \varphi_i(e) = ag mox <math>\theta_i(x,e) is well-defined.
  For (i), clearly I is compact & somex.
  Need to Show that (pls) is continuous (=) Pi(s) is continuous
 Maximum Theorem: Let X,Y S R be compact.
            f: X x Y -> R is continuous & strictly concave in X,
             and \varphi: Y \to X be defined as \varphi(y) = \arg\max_{x \in X} f(x, y)
             Then q(\cdot) is Continuous.
 To show qi (o) is continuous, note that
                     qi(o) = arg max Di(n, o)
ne Zi
                   we've shown D; (n, r) is continuous & strictly concave in
                  the first argunut, Ii. 5 are compact
 Hence by maximum theorem, \varphi(s) is continuous.
 Thus by BFPT, \varphi: \Sigma \to \Sigma has fixed - pt.
  remme: Let \varphi(\hat{s}) = \hat{s} for s \in \Sigma. Then \hat{s} is a NE.
  Proof: & 2 (&, ..., &). & of uniquely maximized
                   \Theta_i(n,\hat{\sigma}) = u_i(n,\hat{\tau}_i) - \|n - \hat{\sigma}_i\|_2^2
                  Say à is not a Nf. Thus Ji, oi E Zi s.t.
                                u_i(\sigma_i', \hat{\sigma_i}) > u_i(\hat{\sigma_i}, \hat{\sigma_i})
                  Let \sigma_i^{\lambda} = \lambda \sigma_i^{i} + (1-\lambda) \sigma_i^{i}
                  \forall \forall \exists \lambda \in (o, i] \text{ s.t. } \theta_i(\sigma_i^{\lambda}, \hat{\sigma}) > \theta_i(\hat{\sigma}_i^{\lambda}, \hat{\sigma})
                            for a contradiction.
                  Now, \theta_i(\delta_i^{\lambda}, \delta) - \theta_i(\delta_i^{\lambda}, \delta)
                                    = u_i(\hat{c_i}^{\lambda}, \hat{c_i}) - u_i(\hat{c_i}, \hat{c_i}) - \left( \|\hat{c_i}^{\lambda} - \hat{c_i}\|_{2^{\perp}} \right)
                                    =\lambda\left(u_{i}\left(\mathscr{E}_{i}^{\prime},\,\,\widehat{\mathscr{E}}_{i}\right)-u_{i}\left(\widehat{\mathscr{E}}_{i},\,\,\widehat{\mathscr{E}}_{-i}\right)\right)
  (skipping some
                                              - \lambda^2 \parallel \sigma_i' - \hat{\sigma_i} \parallel_2^L
       Calculations)
                     for \lambda \in \left(0, \frac{u((\varepsilon', \dot{\varepsilon}))}{\|\varepsilon' - \hat{\varepsilon}\|_{L^{2}}^{2}}\right), this is > 0
               hence, \exists \lambda > 0 s.t. \theta_i(\sigma_i^{\lambda}, \hat{\sigma}) > \theta_i(\hat{\tau}, \hat{\sigma})
                 & hence & is not maximizer of \theta_i(\cdot, 2), giving a
                Contra di ctron.
        This completes proof of Nash Theorem.
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